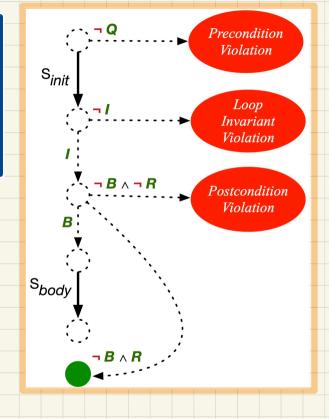
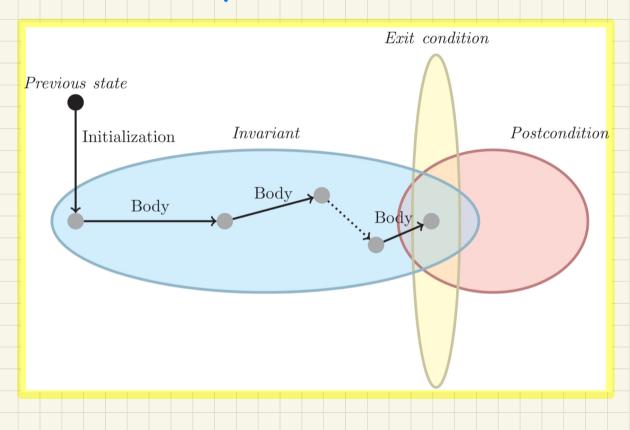
programming statements. Correctness of Loops: Syntax Precondition Precondition void myAlgorithm() { **Violation** assert Q; /* Precondition */ S_{init} S_{init} while (B) assert I; /* Is LI established? */ while(B) { Loop Spody S_{body} Invariant assert I; /* Is LI preserved? */ **Violation** {**R**} assert R; /* Postcondition */ postandippin. $\neg B \land \neg R$ **Postcondition Violation** is B: stay condition
7 B: exit condition. S_{body} * Initialization/Reparation for therations. 4x Ax long as B is true, execute I rody another Ax soon as B is take, exit from the bop. $\neg B \land R$

Correctness of Loops: Example

```
void testLI() { /* Assume: integer attribute i */
assert i == 1; /* Precondition */
assert (1 <= i) && (i <= 6); /* Is LI established? */
while (i <= 5) {
    i = i + 1;
    assert (1 <= i) && (i <= 6); /* Is LI maintained? */
}
assert i == 6; /* Postcondition */
}</pre>
```



Contracts of Loops: Visualization



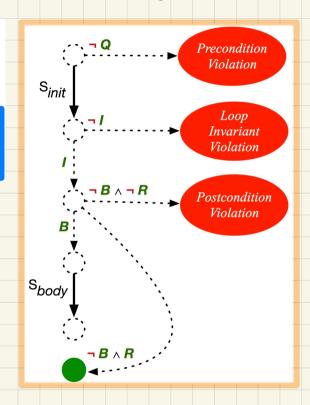
Correctness of Loops: Dijkstra's Shortest-Path Algorithm

Recall: A *loop invariant* (*LI*) is a Boolean condition.

- LI is establised before the 1st iteration.
- *LI* is <u>preserved</u> at the end of each subsequent iteration.

The (iterative) Dijkstra's algorithm has LI:

For every vertext u that has already been $\underline{removed}$ from the priority queue Q (i.e., u is considered $\underline{visited}$), D(u) equals the \underline{true} shortest-path distance from source s to u.



Dijkstra's Shortest Path Algorithm: Negative Weights

The (iterative) Dijkstra's algorithm has LI:

For every vertext u that has already been $\underline{removed}$ from the priority queue Q (i.e., u is considered $\underline{visited}$), D(u) equals the \underline{true} shortest-path distance from source s to u.

